

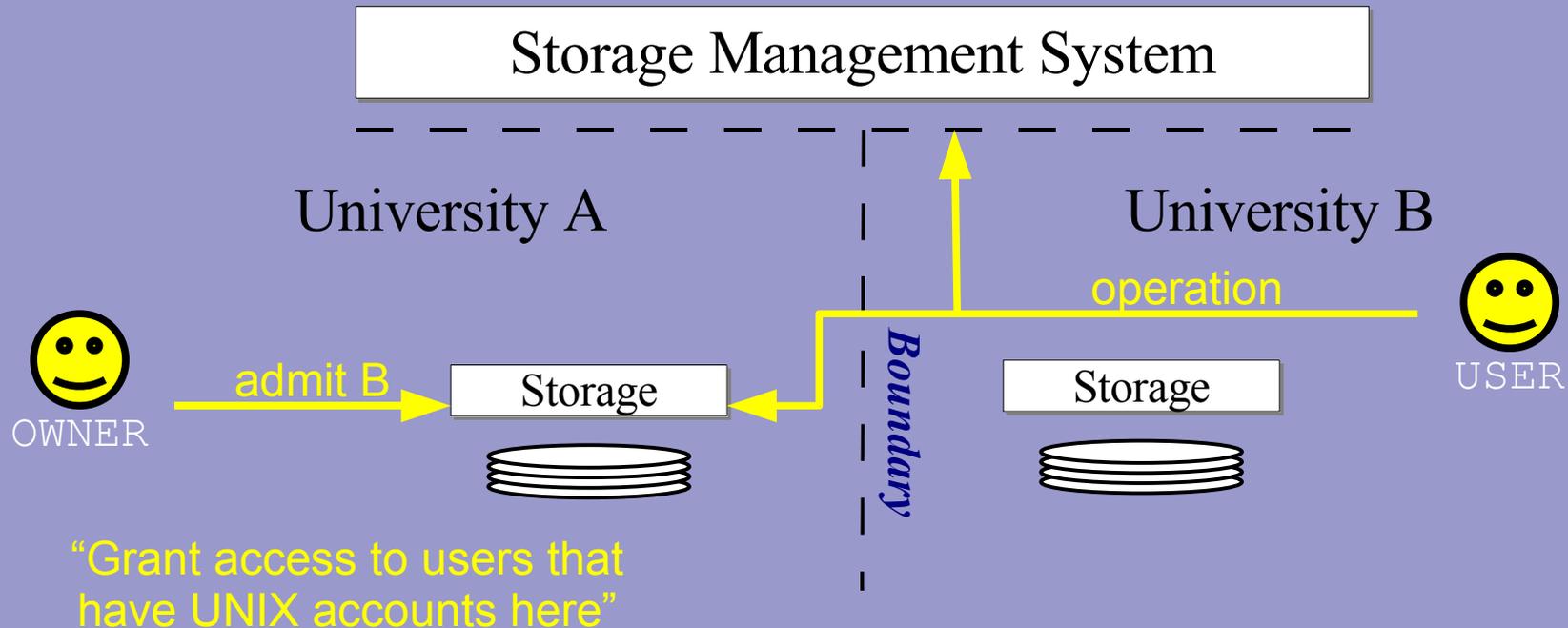
# Access Control for a Replica Management Database

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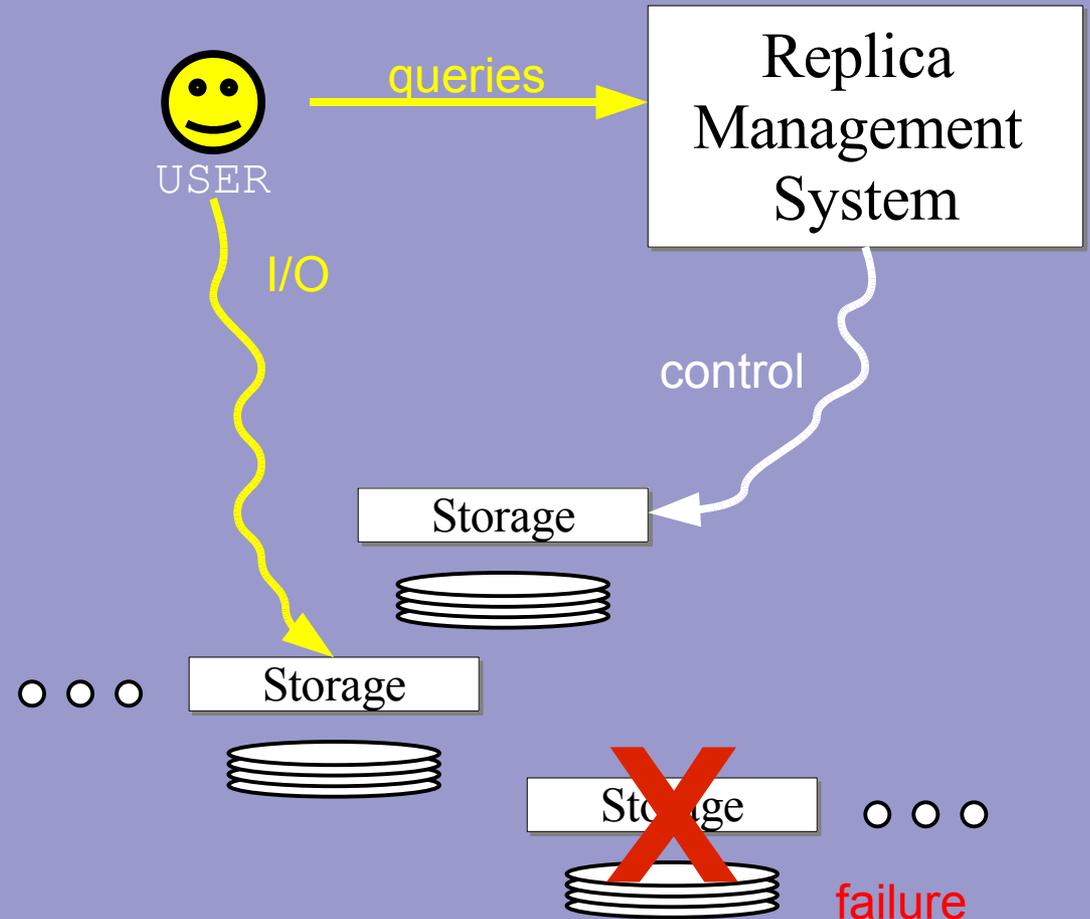
# Ad Hoc Access Control



- Storage owners can grant ad hoc access for remote users
- Such usage may force a change in an overarching system
- How do we make authenticate these operations - without centralized administration?

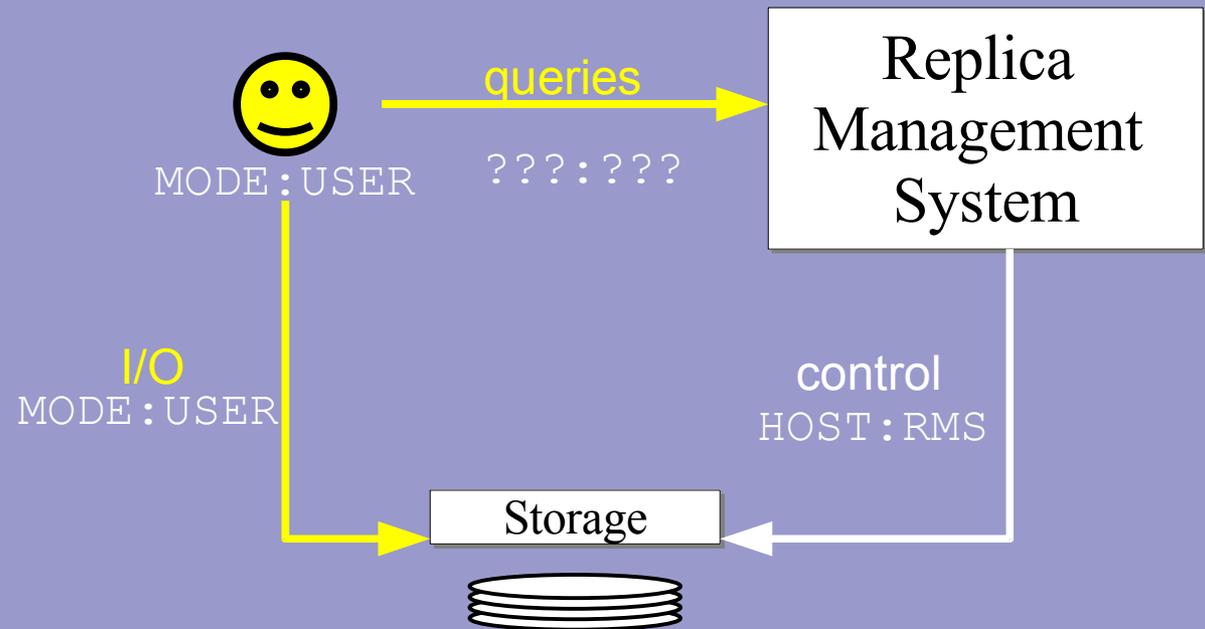
# Storage Fabric

- Cooperative storage clusters in collaborative research environments
- Highly volatile connections
- Multiple-use machines
- Varying free disk space, CPU load, etc.
- *No master list of storage resources...*

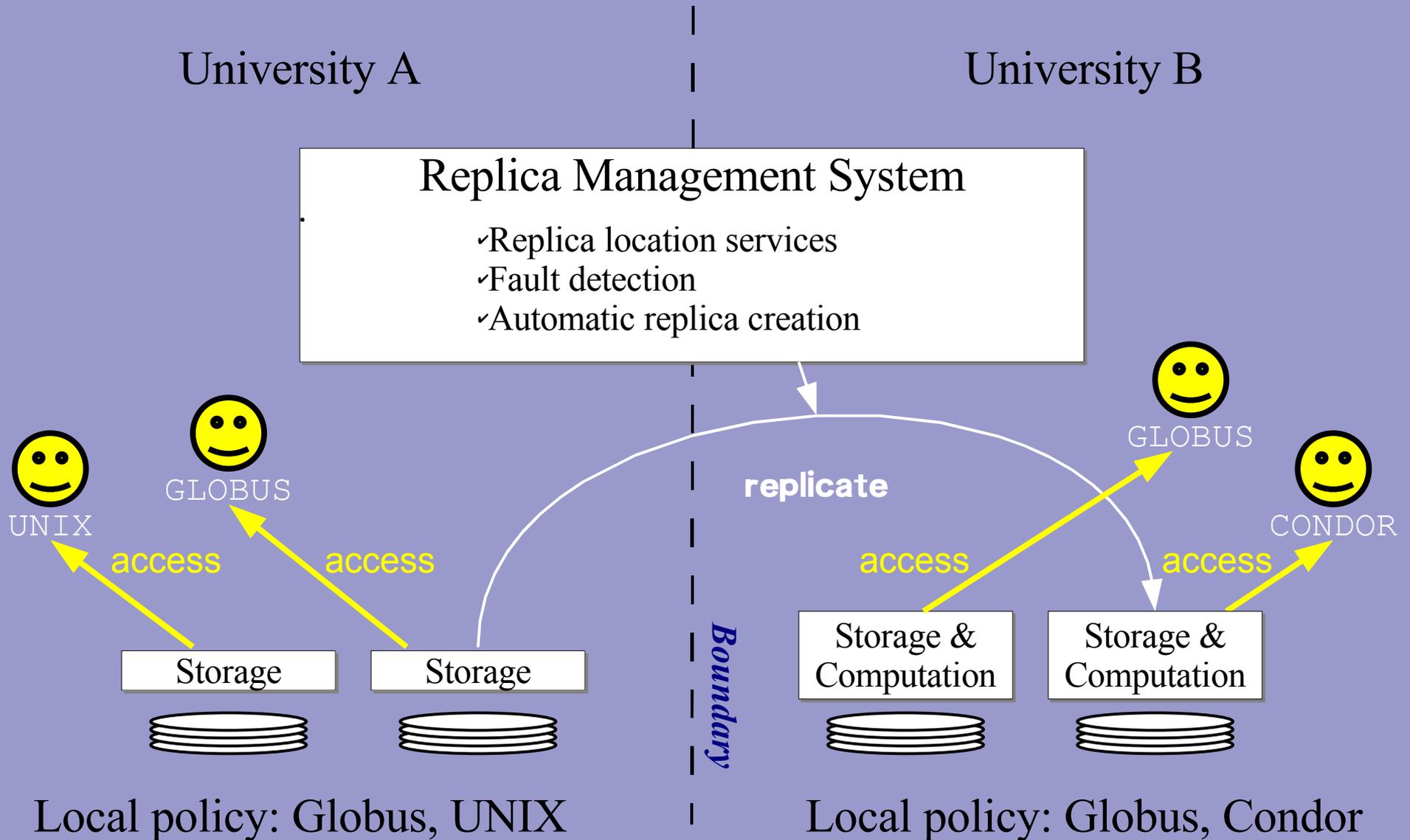


# Plethora of Access Modes

- Cooperative storage users in collaborative research environments
- Highly volatile user groups
- Multiple-identity users
- Varying access to subsets of available systems
- *No master list of users...*



# Ad Hoc Storage Networks



# Outline

- Data description
- The GEMS architecture
- Protocol explanation
- Usage scenarios
- Conclusion

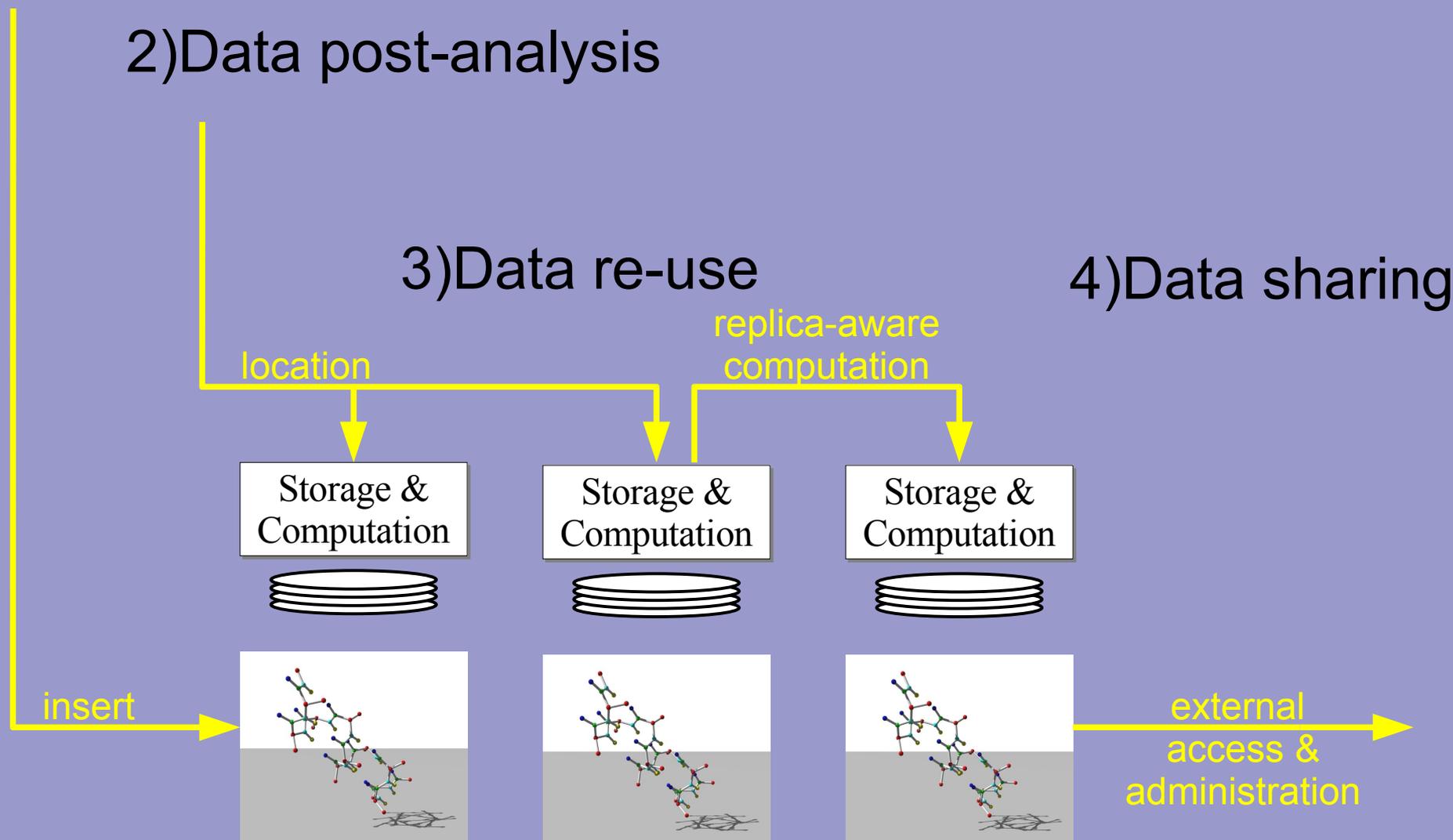
# Grid Enabled Molecular Simulation

1) Data creation

2) Data post-analysis

3) Data re-use

4) Data sharing



# Data Properties

- Start with large data sets in directory trees
- Management aware of data properties
- Records define where they may be stored, who may access them
- Fields specify a management plan

## Replica Management Database

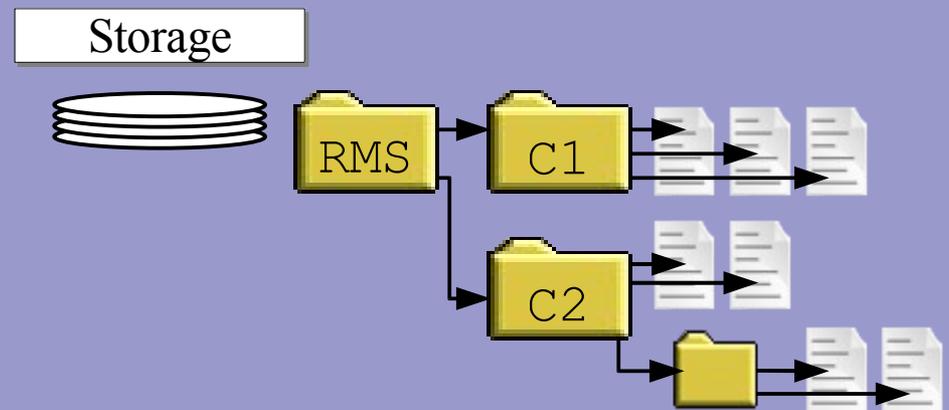
For config  $C_i$ :

METADATA:  $k_1=v_1, \dots$

FILES:  $\{f_1[3] @ \{host_1, host_2, \dots\} \dots\}$

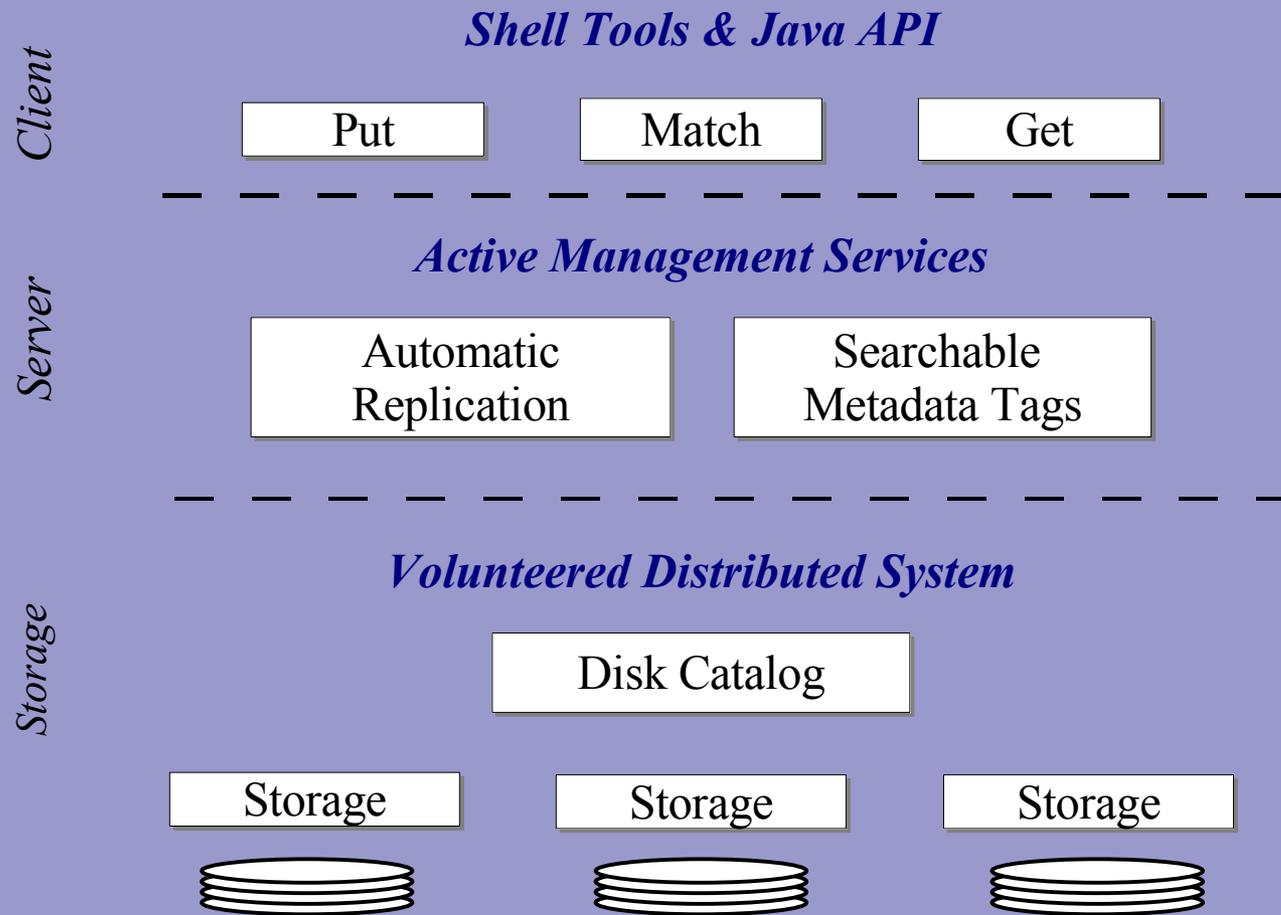
MAP:  $\{cluster_1 \text{ is } \{host_1\},$   
 $cluster_2 \text{ is } \{host_5, host_7\} \dots\}$

ACL:  $\{HOST:RMS \text{ has } \mathbf{RA},$   
 $UNIX:USER1 \text{ has } \mathbf{RA},$   
 $UNIX:USER2 \text{ has } \mathbf{R}\}$





# Architecture

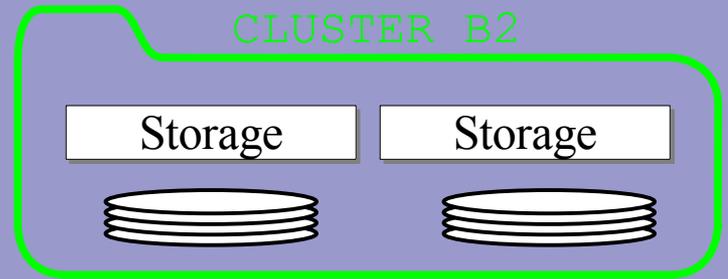
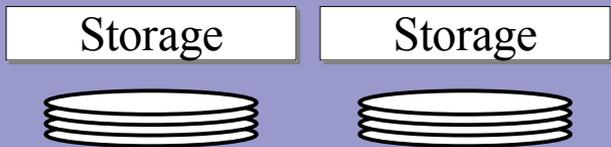
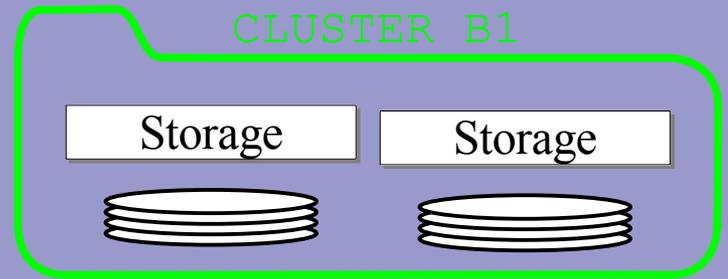
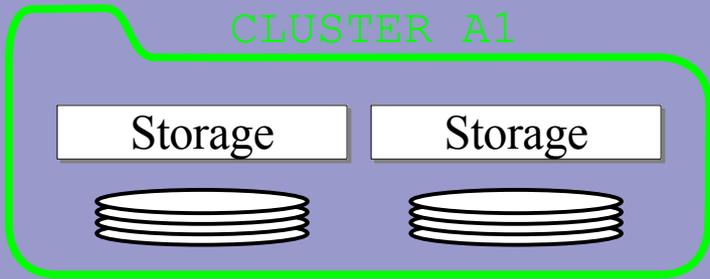


# Record-wise Cluster Topology

University A

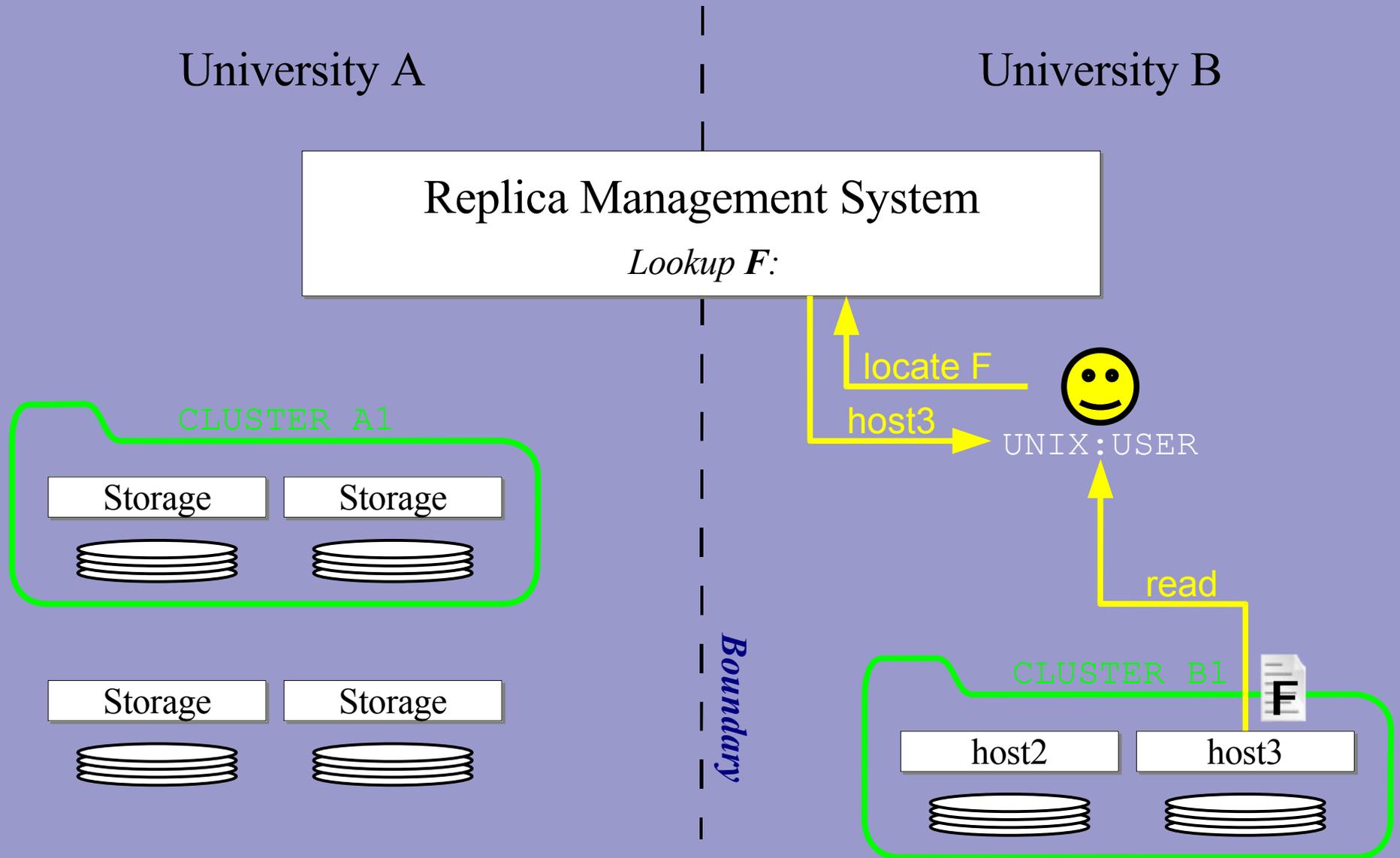
University B

Replica Management System  
*For config  $C_i$ :*



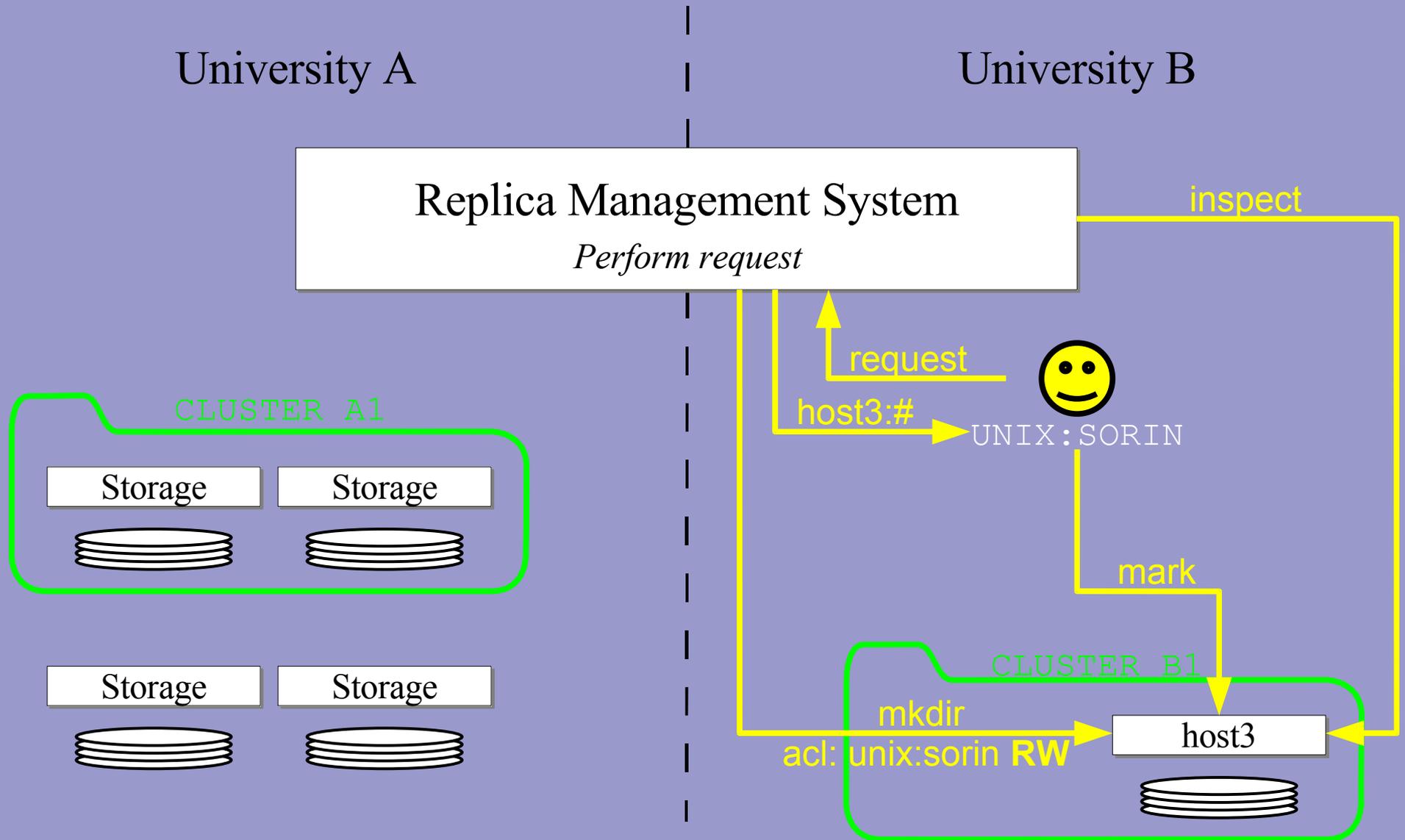
Boundary

# Replica Location and Access



- Topology considered when locating replicas

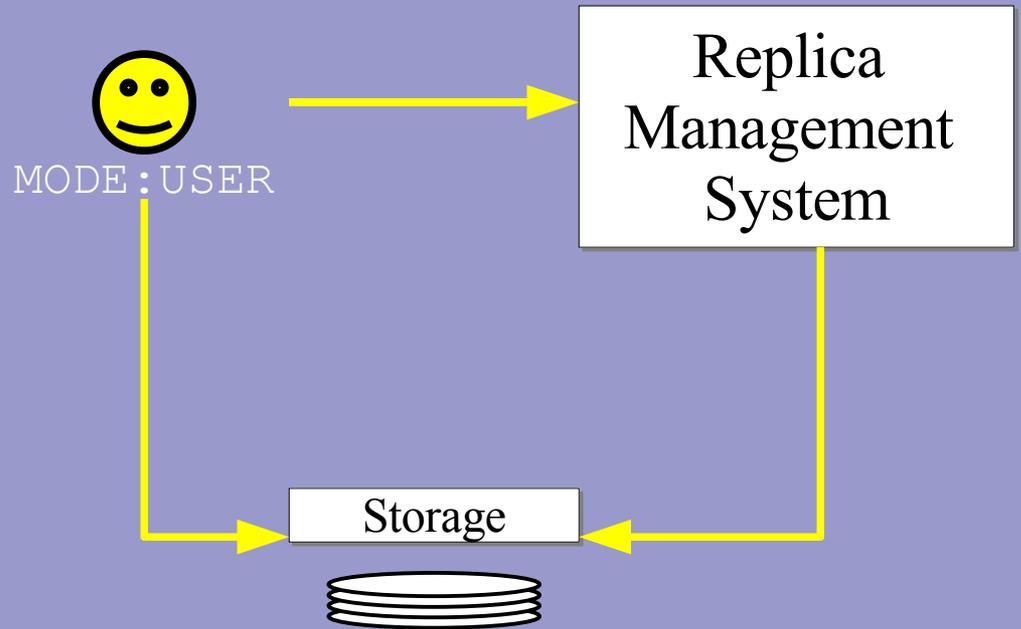
# Database Access



- Database affected by config deletion, modification

# Channel Authentication

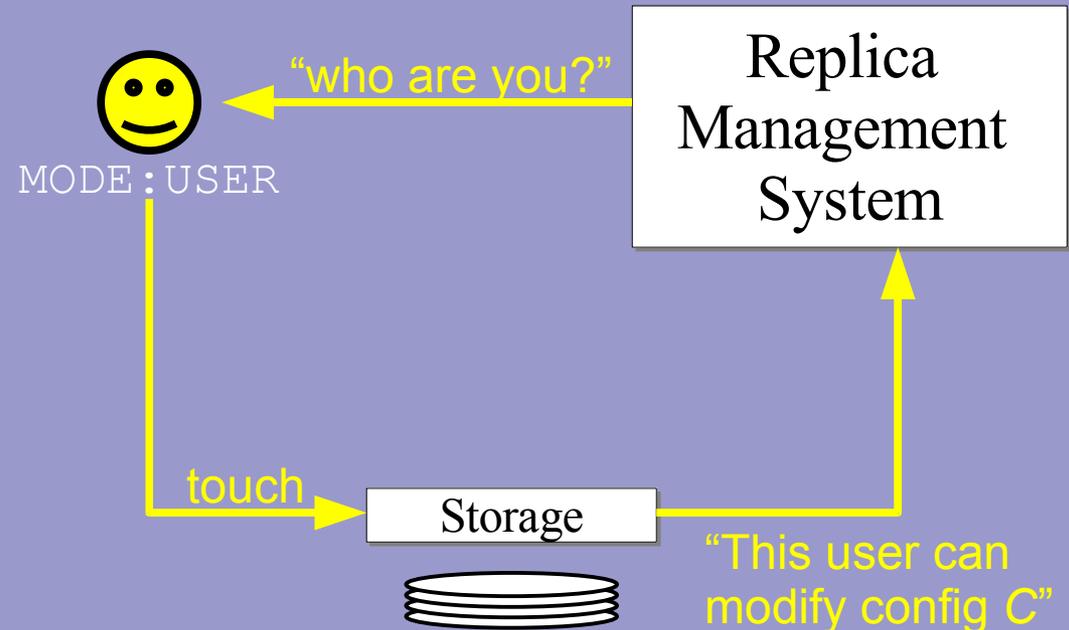
- RMS must determine the ability of a channel to affect a given record
- Record is dependent on storage servers
- RMS may rely on storage servers to “speak for” users
- Users specify permitted storage servers via per-record storage map



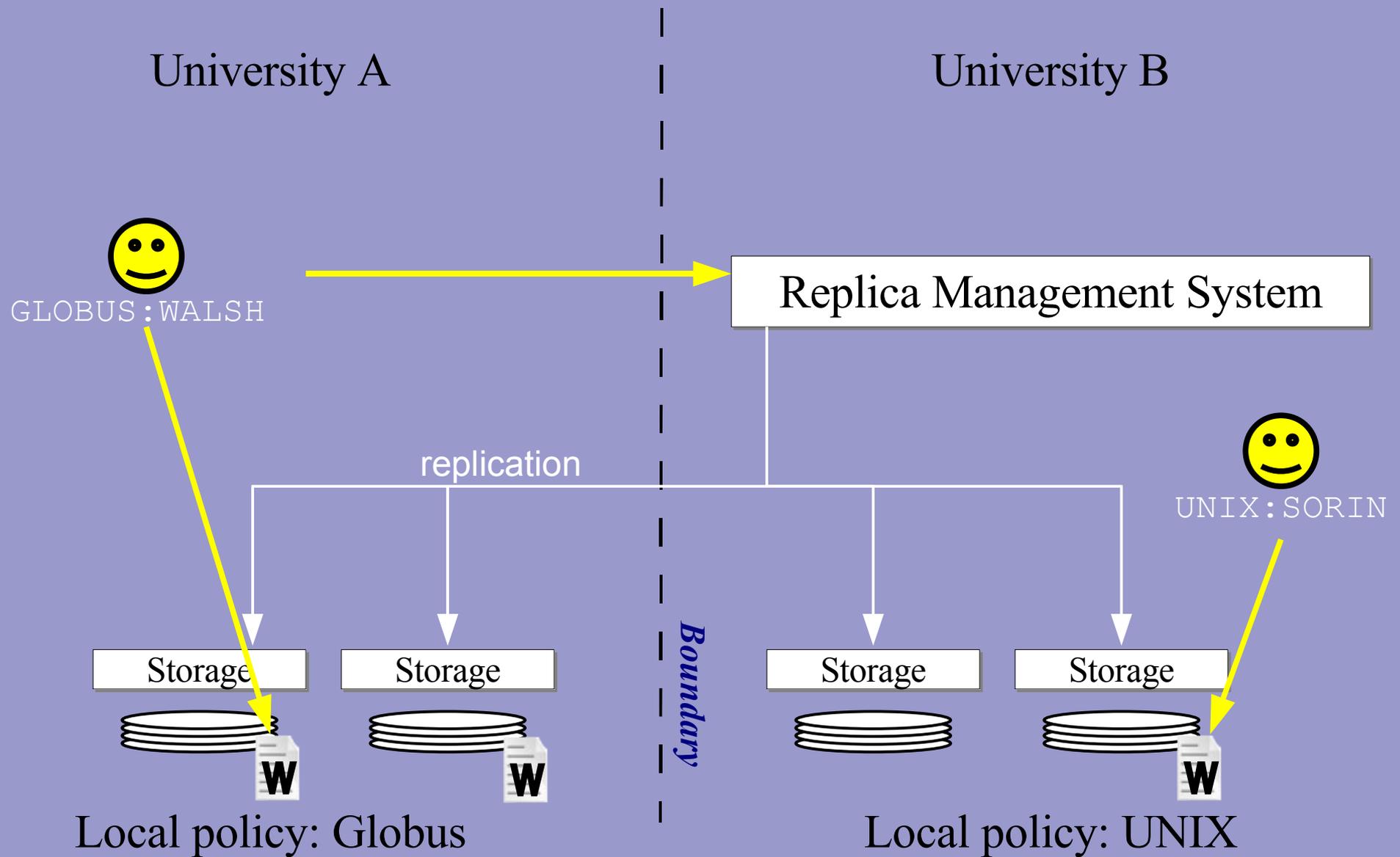
- Mode, user are opaque and are treated as symbols

# Rendition Protocol

- Client requests challenge for config  $C$  on channel  $N$
- RMS creates challenge using eligible host and ACL for  $C$
- Client meets challenge
- RMS checks challenge
- Channel is authenticated for config  $C$

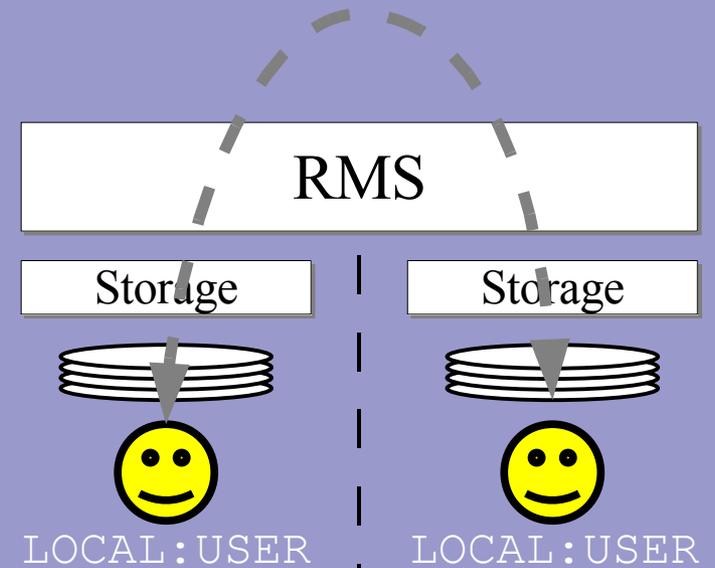


# Cross-boundary Scenarios



# Highlights

- RMS treats users as symbols
- Users may administer data on machines to which they cannot directly authenticate
- Remote replicas are very useful: running remote jobs, collaboration
- Data storage system is essentially as secure as the data creation system or storage servers
- Secure replica placement becomes crucial



- New collaboration technique built around a replica system

# Conclusion

- Heavy emphasis on widely distributed, uncontrolled nature of storage fabric.
- Able to authenticate users indirectly via rendition protocol.



- Generosity and Gluttony in GEMS: Grid Enabled Molecular Simulations.  
Justin Wozniak, Paul Brenner, Douglas Thain, Aaron Striegel, Jesus Izaguirre  
In Proceedings of IEEE High Performance Distributed Computing, July 2005.
- *GEMS Home*: <http://gipse.cse.nd.edu/GEMS>